

# Quantisation feasibility and performance of RSS-based secret key extraction in VANETs

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# Outline

Vehicle Ad-hoc NETWORKS

Physical Layer Security (PLS)

Aims and objectives

Evaluation metrics

The theoretical model

Simulations

Conclusions

# Vehicle Ad-hoc NETWORKs

- ▶ Decentralised **networks of vehicles' on-board-units** (OBUs) and **road-side units** (RSUs)
- ▶ Provide **Safety-related** services, **Navigation-related** information and **Infotainment**
- ▶ System security challenged by the network properties and service constraints

# VANETs Security

- ▶ **Symmetric** techniques are not applicable
- ▶ Proposed **Public Key Cryptography** has drawbacks

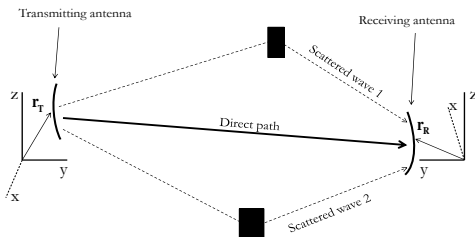
Symmetric (bits)	RSA and Diffie-Hellman (bits)	Elliptic Curve (bits)
80	1024	160
112	2048	224
128	3072	256
192	7680	384
256	15360	512

NIST Recommended Key Sizes

- ▶ Physical layer imperfections can be harnessed to **unconditionally** secure wireless communications

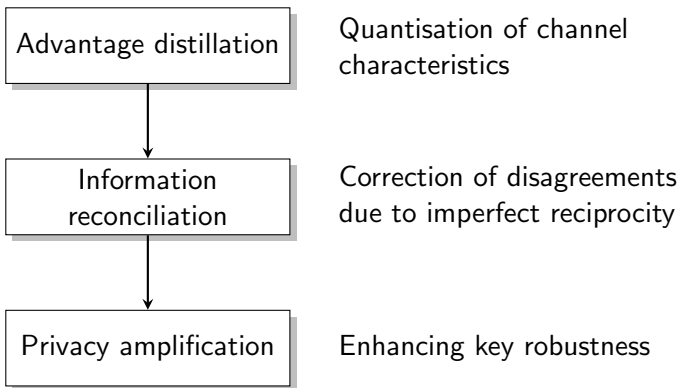
# Physical Layer Security

## Multipath propagation phenomena



- ▶ **Shared** → multipath effects almost identical for both communicating parties (**Channel reciprocity**)
- ▶ **Secret** → correlation rapidly vanishes with time and distance (**Channel variability**)

# The Extraction Process



# Aims and objectives

## Aims

- ▶ Evaluation of RSS-based **Level Crossing** quantisation in VANETs

$$Q(x) = \begin{cases} 1, & \text{if } x > q_+ \\ 0, & \text{if } x < q_- \\ \text{dropped} & \text{otherwise} \end{cases}$$

## Objectives

- ▶ Optimization of system parameters
- ▶ Reliability of the protocol against VANETs requirements
- ▶ Search of possible improvement directions

# Evaluation metrics

- ▶ **Key entropy**

$$H = \sum_{i=0}^N -p_{0,i} \log p_{0,i} - (1 - p_{0,i}) \log(1 - p_{0,i})$$

where  $p_{0,i}$  is the probability of bit  $i$  being 0.

- ▶ **Bit Mismatch Rate** → the ratio of mismatch bits between legitimate parties to the total number of quantised bits
- ▶ **Bit Generation Rate** → the number of secret bits generated per unit time or per sample



## The theoretical model

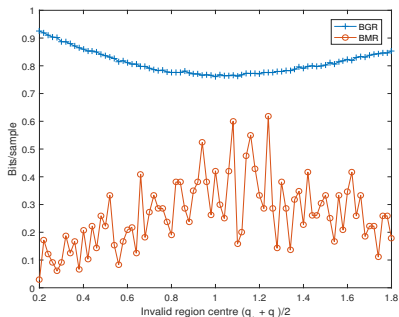
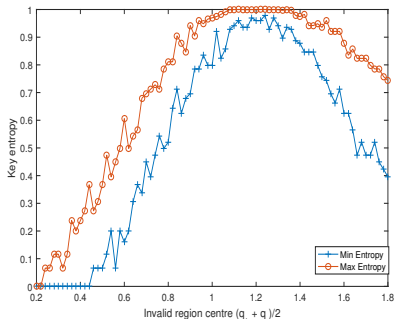
- ▶ Channel modeled as narrow-band frequency-invariant V-V channel with three-dimensional scattering

$$G_N(t) = \sum_{l=1}^L |\alpha_l| \exp(j\phi_l) \exp(j2\pi v_l t)$$

- ▶  $L = 20$  multipath components with constant magnitude  $|\alpha_l|$  and random phase  $\phi_l \sim U[-\pi, \pi]$
- ▶ the Doppler effect  $v_l$  adds up the contributions of transmitter, receiver and scatterers in a three-dimensional environment

# The simulation 1/2

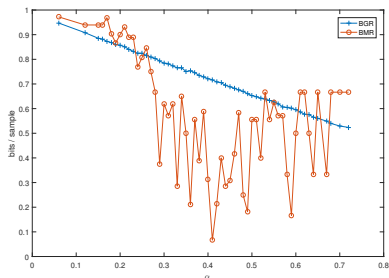
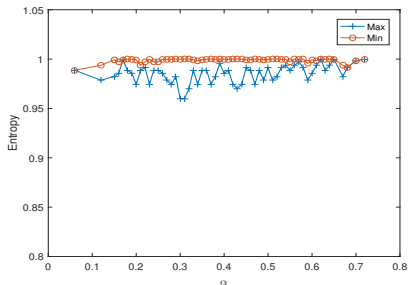
Fixed invalid region size of  $q_+ - q_- = 0.4$



- ▶ Original LC algorithm achieves  $BGR \sim 0.2$  bits/sample, whilst the improved scheme records  $\sim 0.76$  bits/sample
- ▶ Higher results ( $> 0.8$ ) achieved, sacrificing key robustness

## The simulation 2/2

Thresholds computed as  $q_{\pm} = \text{mean}(\hat{h}) \pm \alpha \cdot \text{stdev}(\hat{h})$   
 Parameter  $\alpha$  determines the invalid region size.



- ▶ Optimal value  $\alpha = 0.3$  achieves  $BGR \sim 0.85$  bits / sample

## Conclusions and future work

- ▶ Original LC algorithm generates a shared secret key in not less than two seconds, whilst improved version takes half a second
- ▶ Performances are inadequate if the constraints of safety-related applications (10Hz frequency and  $< 100ms$  latency) are considered
- ▶ Possible improvements
  - ▶ probing rate adaptation
  - ▶ design of a smarter thresholds strategy to exploit the dynamic characteristics of the channel

# Thank you